

Formosa Retreat Juli 2023

2023-07-11



Rosenpass

Wanja Zaeske, Stephan Ajuvo, Marei Peischl,
Benjamin Lipp, Lisa Schmidt, Karolin Varner

<https://rosenpass.eu>

Hello, I am Karolin Varner



- Worked with about every industry tech; incl. Java Web Apps, Microcontrollers, and legacy database system from the 80s
- Did a lot of project management and some people management
- Did a lot of open-source development, privacy- and internet politics activism
- Planning to get involved in the Formosa space

Rosenpass



WireGuard

- ✓ Session-key secrecy
- ✓ ...
- ✓ Identity Hiding
- ✗ Non-Interruptability ¹
- ✗ Post-Quantum Security

PQ WireGuard ²

- ✓ Post-Quantum Security
- ✗ Hybrid security
- ✗ Non-Interruptability ³

Rosenpass

- ✓ Non-Interruptability ⁴
- ✓ Hybrid security ⁵

¹ Assuming a trusted system time

² Hülsing, Ning, Schwabe, Weber, Zimmermann. "Post-quantum WireGuard". <https://ia.cr/2020/379>

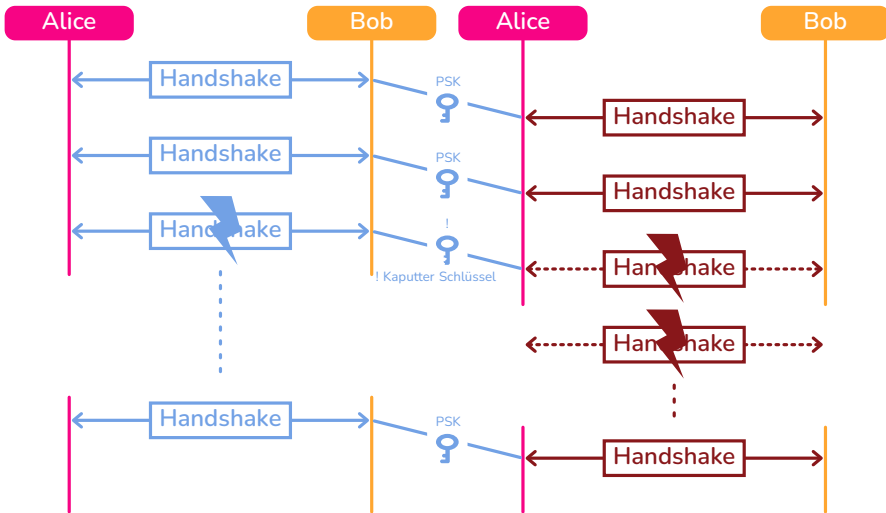
³ Assuming a PSK

⁴ Through cookies

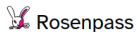
⁵ Used together with standard WireGuard

Rosenpass

WireGuard



Rosenpass can be used right now

[Getting started](#)[About](#)[Source Code](#)[Whitepaper](#)[Contributors](#)[Press](#)[Contact](#)

```
rp pubkey server.rosenpass-secret server.rosenpass-public
rp pubkey client.rosenpass-secret client.rosenpass-public
```

Copy the `-public` directories to the other peers and then start the VPN. On the server:

```
sudo rp exchange server.rosenpass-secret dev rosenpass0 listen 192.168.0.1:9999 \
peer client.rosenpass-public allowed-ips fe80::/64
```

On the client:

```
sudo rp exchange client.rosenpass-secret dev rosenpass0 \
peer server.rosenpass-public endpoint 192.168.0.1:9999 allowed-ips fe80::/64
```

Assign IP addresses:

```
sudo ip a add fe80::1/64 dev rosenpass0 # Server
sudo ip a add fe80::2/64 dev rosenpass0 # Client
```

Test the connection by pinging the server on the client machine:

```
ping fe80::1%rosenpass0 # Client
```

You can watch how Rosenpass replaces the WireGuard PSK with the following command:

```
watch -n 0.2 'wg show all; wg show all preshared-keys'
```

ProVerif in Technicolor



```
~/p/rosenpass ▶ dev/karo/rwpqc-slides ? nix build .#packages.x86_64-linux.proof-proverif --print-build-logs [17/17]
rosenpass-proverif-proof> unpacking sources
rosenpass-proverif-proof> unpacking source archive /nix/store/cznyv4ibwlbzh257v6lzx8r8a14cb0v0-source
rosenpass-proverif-proof> source root is source
rosenpass-proverif-proof> patching sources
rosenpass-proverif-proof> configuring
rosenpass-proverif-proof> no configure script, doing nothing
rosenpass-proverif-proof> building
rosenpass-proverif-proof> no Makefile, doing nothing
rosenpass-proverif-proof> installing
rosenpass-proverif-proof> $ metaverif analysis/01_secrecy.entry.mpv -color -html /nix/store/gidm68r04lkpanvkgz48527qf6nym6dv
-rosenpass-proverif-proof
rosenpass-proverif-proof> $ metaverif analysis/02_availability.entry.mpv -color -html /nix/store/gidm68r04lkpanvkgz48527qf6n
ym6dv-rosenpass-proverif-proof
rosenpass-proverif-proof> $ wait -f 34
rosenpass-proverif-proof> $ cpp -P -I/build/source/analysis analysis/01_secrecy.entry.mpv -o target/proverif/01_secrecy.ent
ry.i.pv
rosenpass-proverif-proof> $ cpp -P -I/build/source/analysis analysis/02_availability.entry.mpv -o target/proverif/02_availab
ility.entry.i.pv
rosenpass-proverif-proof> $ awk -f marzipan/marzipan.awk target/proverif/01_secrecy.entry.i.pv
rosenpass-proverif-proof> $ awk -f marzipan/marzipan.awk target/proverif/02_availability.entry.i.pv
rosenpass-proverif-proof> 4s ✓ state coherence, initiator: Initiator accepting a RespHello message implies they also generat
ed the associated InitHello message
rosenpass-proverif-proof> 35s ✓ state coherence, responder: Responder accepting an InitConf message implies they also genera
ted the associated RespHello message
rosenpass-proverif-proof> 0s ✓ secrecy: Adv can not learn shared secret key
rosenpass-proverif-proof> 0s ✓ secrecy: There is no way for an attacker to learn a trusted kem secret key
rosenpass-proverif-proof> 0s ✓ secrecy: The adversary can learn a trusted kem pk only by using the reveal oracle
rosenpass-proverif-proof> 0s ✓ secrecy: Attacker knowledge of a shared key implies the key is not trusted
rosenpass-proverif-proof> 31s ✓ secrecy: Attacker knowledge of a kem sk implies the key is not trusted
```

Having worked in industry has some advantages



- Knowing how to get projects done
- Coordinating teams instead of working on my own
- Product and user focused perspective
- Building tools that I can use to be more productive
- Open-Source approach: How to catch new contributors



The spec makes it easy to implement Rosenpass

Initiator Code Responder Code Comments

1 InitHello { sidi, epki, sctr, pidiC, auth } 2

Line	Variables ← Action	Variables ← Action	Line
IHR1	<code>ck ← lhash("chaining key init", spkr)</code>	<code>ck ← lhash("chaining key init", spkr)</code>	IHR1
IHR2	<code>sidi ← random_session_id();</code>		
IHR3	<code>eski, epki ← EKEM.keygen();</code>		
IHR4	<code>mix(sidi, epki);</code>	<code>mix(sidi, epki)</code>	IHR4
IHR5	<code>sctr ← encaps_and_mix<SKEM>{spkr};</code>	<code>decaps_and_mix<SKEM>{sctr, spkr, ct1}</code>	IHR5
IHR6	<code>pidiC ← encrypt_and_mix{pidi};</code>	<code>spki, psk ← lookup_peer{decrypt_and_mix{pidiC}}</code>	IHR6
IHR7	<code>mix{spki, psk};</code>	<code>mix{spki, psk};</code>	IHR7
IHR8	<code>auth ← encrypt_and_mix{empty()};</code>	<code>decrypt_and_mix{auth}</code>	IHR8

- Comment**
- Initialize the chaining key, and bind to the responder's public key.
 - The session ID is used to associate packets with the handshake state.
 - Generate fresh ephemeral keys, for forward secrecy.
 - InitHello includes sidi and epki as part of the protocol transcript, and so we mix them into the chaining key to prevent tampering.
 - Key encapsulation using the responder's public key. Mixes public key, shared secret, and ephemeral into the chaining key, and authenticates the responder.
 - Tell the responder who the initiator is by transmitting the peer ID.
 - Ensure the responder has the correct view on spki. Mix in the PSK as optional static symmetric key, with epki and spkr serving as nonces.
 - Add a message authentication code to ensure both participants agree on the session state and protocol transcript at this point.

4 RespHello { sidr, sidi, ecti, scti, biscuit, auth } 3

Line	Variables ← Action	Variables ← Action	Line
RHR1		<code>sidr ← random_session_id();</code>	RHR1
RHR2	<code>ck ← lookup_session{sidi};</code>		RHR2
RHR3	<code>mix{sidr, sidi};</code>	<code>mix{sidr, sidi};</code>	RHR3
RHR4	<code>ecti ← encaps_and_mix<EKEM>{epki, ecti};</code>	<code>ecti ← encaps_and_mix<EKEM>{epki};</code>	RHR4
RHR5	<code>decaps_and_mix<SKEM>{scti, spki, scti};</code>	<code>scti ← encaps_and_mix<SKEM>{spki};</code>	RHR5
RHR6	<code>mix{biscuit};</code>	<code>biscuit ← store_biscuit();</code>	RHR6
RHR7	<code>decrypt_and_mix{auth}</code>	<code>auth ← encrypt_and_mix{empty()};</code>	RHR7

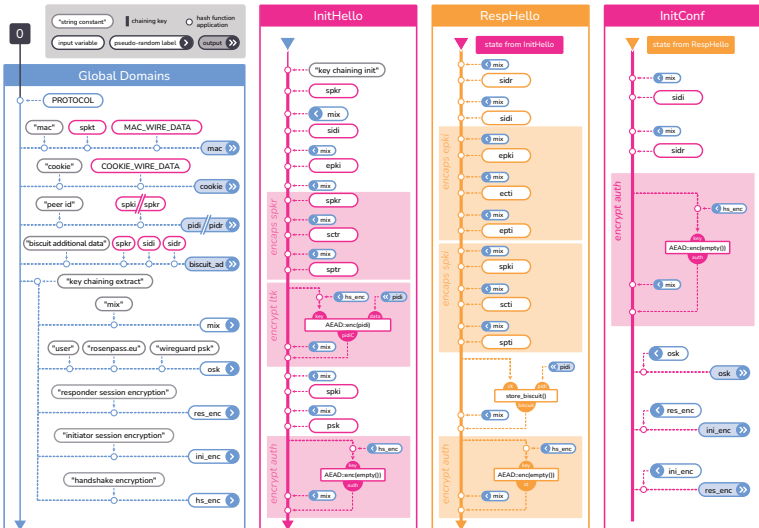
- Comment**
- Responder generates a session ID.
 - Initiator looks up their session state using the session ID they generated.
 - Mix both session IDs as part of the protocol transcript.
 - Key encapsulation using the ephemeral key, to provide forward secrecy.
 - Key encapsulation using the initiator's static key, to authenticate the initiator, and non-forward-secret confidentiality.
 - The responder transmits their state to the initiator in an encrypted container to avoid having to store state.
 - Add a message authentication code for the same reason as above.

5 InitConf { sidi, sidr, biscuit, auth } 6

Line	Variables ← Action	Variables ← Action	Line
ICR1		<code>biscuit_no ← load_biscuit{biscuit};</code>	ICR1
ICR2		<code>encrypt_and_mix{empty()};</code>	ICR2
ICR3	<code>mix{sidr, sidi};</code>	<code>mix{sidr, sidi};</code>	ICR3
ICR4	<code>auth ← encrypt_and_mix{empty};</code>	<code>decrypt_and_mix{auth};</code>	ICR4
ICR5		<code>assert{biscuit_no > biscuit_used};</code>	ICR5
ICR6		<code>biscuit_used ← biscuit_no;</code>	ICR6

- Comment**
- Responder loads their biscuit. This restores the state from after RHR6.
 - Responder recomputes RHR7, since this step was performed after biscuit encoding.
 - Mix both session IDs as part of the protocol transcript.
 - Message authentication code for the same reason as above, which in particular ensures that both participants agree on the final chaining key.
 - Biscuit replay detection.
 - Biscuit replay detection.

Professional illustrators create stunning graphics



Creating successful projects by knowing what not to do



- Rosenpass avoids targeting: GUIs, VPN data transport, support for many platforms
- Instead we: Created a core technology; working with companies to integrate Rosenpass (e.g. Open-Source VPN startups)
- **Vitaly we** chose to focus on API; making it easy to integrate Rosenpass
- **Vitaly we** integrate with the existing ecosystem (i.e. WireGuard) instead of trying to replace it

Starting partnerships...



- with Open-Source VPN companies
- with Kubernetes VPN companies
- with Quantum-Key-Distribution Projects
- to verify the Rosenpass source code
- to apply isolation features to Rosenpass (Micro-VMs)
- with university teaching departments to use the project as a simple example of bleeding-edge modern crypto

Talk to me about...



- using Rosenpass as demonstrator-project to integrate new cryptographic technologies in
- figuring out how to attract independent contributors to Formosa Projects
- applying API-focused techniques to Formosa projects to emphasize interoperability
- **Idea:** Providing XML-representations of proof assistants'⁶ inputs and outputs to allow easy integration with external tools
- **Idea:** Python libraries to work with Formosa tools as everybody knows python

⁶ EasyCrypt, CryptoVerif, ProVerif, Tamarin