

<https://rosenpass.eu>

<https://chaos.social/@rosenpass>

MRMCD 2023

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Sichere Kryptografie trotz Quantencomputern: Projektupdate

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Funding: NLNet & Prototype Fund



Was passiert im Talk?

- Was bisher geschah
- Zusammenfassung vom EH20 Talk: Was ist Rosenpass
- Was nach dem Easterhegg passiert ist
- Was wir nun vor haben
 - go-rosenpass
 - NetBird
 - Broker-Architektur & Schnittstellen zum Einbinden



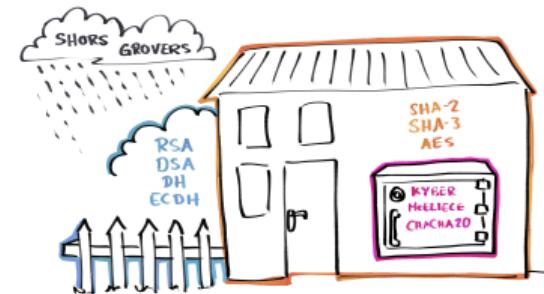
Was bisher geschah

- Seit 2020: Entwicklung der Kryptografie & der Software
- Feb. 2023: Softwarerlease & Whitepaper
- März 2023: NLNet Projekt um Sicherheitsbeweis mit CryptoVerif zu erzeugen
- März 2023: Talk auf dem Real World Post-Quantum Krypto Workshop in Tokyo
- April 2023: Vorstellung/Erklärung auf dem Easterhegg¹
- Aug. 2023: Release Kandidat 0.2.0 mit FreeBSD Unterstützung
- Sep. 2023: Beginn des Prototype Fund 14 Projektes für Isolation in Rosenpass

¹ <https://media.ccc.de/search/?q=rosenpass>

Warum sind Quantencomputer (k)eine Bedrohung?

- Grovers Algorithmus **schwächt** symmetrische Kryptografie
 - AES, SHA-2, SHA-3, Chacha20
 - Lösung: größere Keys
- Shors Algorithmus **bricht** asymmetrische Kryptografie
 - RSA, DSA, DH, ECDH
 - Lösung: alternative Kryptografie
- Nur auf großen Quantencomputern
 - Die existieren noch nicht
 - Problem: Store now, decrypt later



Quantencomputer überschatten Kryptografieverfahren.



PQ-sichere VPNs: WireGuard + Rosenpass

- Hybride Sicherheit
 - Bricht nur, wenn Rosenpass **und** WireGuard versagen
- Überall nutzbar, wo WireGuard schon läuft
- Ohne Anpassung vom WireGuard Source Code
 - Shared Secret aus Rosenpass = PSK für WireGuard
- Aber:
 - Ein Prozess mehr
 - Handshake alle 2 Minuten



WireGuard mit Rosenpass.



Rosenpass: Sicherheitseigenschaften

WireGuard

PQ WireGuard³

Rosenpass

Session-key secrecy

Post-Quantum Security

Non-Interruptability⁵

...

Hybrid security

Hybrid security⁶

Identity Hiding

Non-Interruptability⁴

Non-Interruptability²

Post-Quantum Security

² Angenommen der Systemzeit wird Vertraut

³ Hülsing, Ning, Schwabe, Weber, Zimmermann. "Post-quantum WireGuard". <https://ia.cr/2020/379>

⁴ Assuming a PSK

⁵ Through cookies

⁶ Wenn es mit WireGuard benutzt wird



Zum Nachbauen... aus dem Whitepaper:

Initiator Code	Responder Code	Comments																																			
1 InitHello { sidi, epki, sctr, pidiC, auth }	2																																				
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Zum Nachbauen... go-rosenpass – Steffen Vogel FTW

Screenshot of the GitHub repository page for `cunicu/go-rosenpass`.

The repository is public and has 12 issues, 1 pull request, and 3 watchers.

The main branch is `main`. The `About` tab is selected.

A port of Rosenpass post-quantum key-exchange protocol to Go.

Tags: `go`, `golang`, `cryptography`, `vpn`, `wireguard`, `post-quantum`, `rosenpass`.

Recent commits:

- `stv0g` Use wg-quick configuration files rather netlink interface to ge... (yesterday) 146
- `.github/workflows` Run CI tests with root permissions (yesterday)
- `.reuse` Update name of Go module and fix links for new re... (3 weeks ago)
- `.vscode` Implement integration / interoperability tests (4 months ago)
- `LICENSES` make repo REUSE compliant (4 months ago)
- `cmd` Add new exchange-intf and gen-keys-intf sub-com... (yesterday)
- `config` Use wg-quick configuration files rather netlink inte... (yesterday)

Repository statistics:

- 4 stars
- 3 watching
- 1 fork

Page footer: 7/1



Zum Nachbauen... go-rosenpass – Steffen Vogel FTW

rosenpass / rosenpass

Code Issues 41 Pull requests 1 Discussions Actions Projects Wiki Security ...

Whitepaper proof read #68

Open 17 tasks stv0g opened this issue on May 20 · 26 comments

stv0g commented on May 20 • edited

I've started to work on a Golang implementation of the Rosenpass key exchange^[1]. While implementing it, I stumbled over some confusing parts in the whitepaper:

Variables / message fields mac, cookie, ini_enc & res_enc

Figure 3 shows the output variables mac and cookie which apparently should be included in a message envelope (Figure 2). But I find no mention about use of the mac and cookie variable elsewhere in the whitepaper.

Apart from the osk two more output variables are generated after the completion of the handshake: ini_enc & res_enc

Assignees: No one—assign yourself

Labels: None yet

Projects: Rosenpass 1.0.0

Status: No status



Zum Nachbauen... go-rosenpass – Steffen Vogel FTW

[Open](#) [17 tasks](#) **Whitepaper proof read #68**
stv0g opened this issue on May 20 · 26 comments

Status
<ul style="list-style-type: none"><input type="checkbox"/> Describe Payload transmission as a possible protocol extension<input type="checkbox"/> Highlight relation to WireGuard whitepaper<input type="checkbox"/> How are the roles of initiator / responder assigned?<input type="checkbox"/> Endianess of biscuit counter<input type="checkbox"/> Figure 4: Wrong cipher-text variable<input type="checkbox"/> Section 2.4.1: Better naming for session/index table<input type="checkbox"/> Inverted assertions for replay detection<input type="checkbox"/> Figure 3: Wrong PRF Labels for chaining key extract/init #67<input type="checkbox"/> Figure 3: Wrong PRF labels for session encryption keys<input type="checkbox"/> Section 2.3: Wrong protocol identifier<input type="checkbox"/> Section 2.1.1: Wrong hash function?<input type="checkbox"/> Section 2.1.1: Describe non-standard HMAC-Blake2 variant<input type="checkbox"/> Figure 3 / Section 2.5: Wrong order of mixing in en/decaps_and_mix()<input type="checkbox"/> Figure 2: Wrong field ordering in RespHello message<input type="checkbox"/> Section 2.4.3: Mismatching biscuit key epoch<input type="checkbox"/> Section 2.1.4 / 2.1.5: Add references to exact versions of KEM specifications which are used by Rosenpass



Zum Integrieren... NetBird

SECURITY.md	Add security policy file (#600)	10 months ago
go.mod	Routemgr error handling (#1073)	3 weeks ago
go.sum	Routemgr error handling (#1073)	3 weeks ago

[README.md](#)

New Release! Self-hosting in under 5 min. [Learn more](#)

netbird

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[slack](#) [@netbird](#)

Start using NetBird at [netbird.io](#)

[See Documentation](#)

[Join our Slack channel](#)

NetBird combines a configuration-free peer-to-peer private network and a centralized access control system in a single platform, making it easy to create secure private networks for your organization or home.

Languages

Go 97.4%	Shell 2.4%
● Other 0.2%	



Zum Integrieren... NetBird

☰ README.md

Together with CISPA Helmholtz Center for Information Security NetBird brings the security best practices and simplicity to private networking.



Testimonials

We use open-source technologies like [WireGuard®](#), [Pion ICE \(WebRTC\)](#), [Coturn](#), and [Rosenpass](#). We very much appreciate the work these guys are doing and we'd greatly appreciate if you could support them in any way (e.g. giving a star or a contribution).

Legal

WireGuard and the *WireGuard* logo are registered trademarks of Jason A. Donenfeld.

Zum Integrieren... NetBird

- Netbird: Einfaches user interface
- Rosenpass: Hochsichere PQ-Crypto
- go-rosenpass: Um Plattformen zu unterstützen auf denen die Rust Variante schwer zu integrieren ist
 - Android
 - iOS
 - Windows
 - ...



Zum Integrieren... Prototypefund 14 Projekt

- Schnittstelle zwischen Komponenten
- Kommunikation über Unix-Sockets
- Spezielle Serialisierungsbibliothek für Schlüsseldaten⁷
- Broker-Pattern für Rosenpass– Jede Komponente in einem eigenen Prozess
- Mikro-VMs um wirklich hohe Sicherheit zu haben
- Minimale Privilegien; Sandboxing



⁷ Externes Memory-Management



Rosenpass Roadmap

- Sicherheitsbeweis
- Formelle Verifikation der Implementierung
- Einfachere Benutzbarkeit
- Kryptografie + Safety Forschung:
 - Kryptografie in der Avionik
 - Decryption Despite Error





Rosenpass Strukturpläne

- Translationsforschung: Schnittstelle zwischen Industrie und Wissenschaft
- Mit mehreren Integratoren arbeiten; Open-Source R&D-Abteilung
- Antihierarchisches Arbeiten
- Karo hätte gerne mal wieder Freizeit



Idee: Kryptografie als Notar Erklären

- Problem: Kryptografie wird als schwarze Magie verstanden
- Problem: Krude Vorschläge zur Verwaltung automatisierung
- Problem: Und Strafverfolgung
- Problem: Krypto wird auf Verschlüsselung reduziert
- Problem: Kaum jemand weiß was moderne Verfahren tun
 - Elliptic-Curve Pairings
 - Multi-Party Computation
 - Homomorphe Verschlüsselung
 - Datenbanken mit anonymem Zugriff
 - Anonyme Kommunikation



Quelle: White Rabbit aus Alice in Wonderland – CC0

Häschen sind die besseren Menschen.

Idee: Kryptografie als Notar Erklären

- Krypto: Einsatzfähig für viele Prozesse in denen Information Übertragen wird
- Datenschutz: Anonyme, Nutzergesteuerte Prozesse
- Idee: Metapher von Kryptografie als Notar
 - Notare werden Bestraft wenn sie Dinge Zusichern die sie nicht können
 - Oder wenn sie gegen Regeln verstößen
 - Spezieller schutz vor dem Recht
 - Kryptografie: Ähnlich, nur Matematisch, statt mit Staatsgewalt



Quelle: White Rabbit aus Alice in Wonderland – CC0

Häßchen sind die besseren Menschen.